Pulse sequences for suppressing leakage in single-qubit gate operations

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Many realizations of solid-state qubits involve couplings to leakage states lying outside the computational subspace, posing a threat to high-fidelity quantum gate operations. Mitigating leakage errors is especially challenging when the coupling strength is unknown, e.g., when it is caused by noise. Here we show that simple pulse sequences can be used to strongly suppress leakage errors for a qubit embedded in a three-level system. As an example, we apply our scheme to the recently proposed charge quadrupole (CQ) qubit for quantum dots. These results provide a solution to a key challenge for fault-tolerant quantum computing with solid-state elements.

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Recent advances in semiconducting quantum dots make them promising candidates for universal quantum computing [1–6]. However, performing gate operations with high enough fidelity to support fault-tolerant error correction remains a key challenge [7]. Two ways that a qubit can fail to have high fidelity are (i) the qubit could decay or dephase, and (ii) quantum information could leak out of the qubit’s logical subspace into other quantum states in the physical system [8–12]. While several recent proposals in quantum dots have focused on suppressing dephasing from environmental noise [13,14], relatively little effort has gone into suppressing leakage [12].

Several approaches for reducing leakage errors have been developed for superconducting qubits, including analytic pulse shaping [9,15] and optimal quantum control [16,17]. It is also known that leakage errors are, in principle, suppressible for a system-bath model with a composite sequence of a large number of pulses [8]. However, suppressing leakage errors below the fault-tolerant threshold under experimental conditions for quantum dot qubits is challenging because of the need to apply smoothly varying short control pulses [18] within a time much less than the coherence times of the system and also because the qubits experience fluctuations that vary in time and strength [19]. For semiconducting systems, these constraints preclude using analytically derived pulse shapes, since these require knowing the noise strength as an input parameter [9], or quantum control strategies that rely on optimizing an unrealistically large number of control parameters [16,17].

In this Rapid Communication, we develop a simple and experimentally feasible protocol based on composite pulses that suppresses leakage errors in a three-level quantum system where the leakage state is coupled to one of the logical states with an unknown but static coupling strength (referred to as the quasistatic noise approximation [19]). Whereas the leakage error scales quadratically with noise amplitude in conventional pulsed-gate schemes for this model, our protocol significantly improves the gate fidelities by eliminating computational errors in the logical subspace up to fourth order in the noise amplitude and leakage errors up to sixth order.

In a basis comprised of two logical states and one leakage state, the model Hamiltonian is given by

\[ H = H_z + H_x + H_{\text{leak}}, \]

with

\[ H_z = \frac{\epsilon_q}{2} \begin{pmatrix} 0 & 0 & 0 \\ 0 & -1 & 0 \\ 0 & 0 & -\zeta \end{pmatrix}, \quad H_x = g \begin{pmatrix} 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 0 \end{pmatrix}, \]

and

\[ H_{\text{leak}} = \xi \begin{pmatrix} 0 & 0 & 0 \\ 0 & 0 & 1 \\ 0 & 1 & 0 \end{pmatrix}. \]

where \( \epsilon_q \) and \( g \) are the independent control parameters for rotations about the \( z \) and \( x \) axes of the Bloch sphere in the logical subspace, \( \xi \) denotes the unknown coupling between the leakage state and one of the logical states, and \( \zeta \) denotes the (scaled) leakage state energy in the absence of coupling. While the Hamiltonian (1) provides a very general description of a two-level system coupled to a leakage state [9,20], we focus here on the semiconducting charge quadrupole (CQ) qubit [21], for which the logical states have different charge and quadrupolar detuning parameters, respectively. It is inherent that the states are protected from the predominant type of noise in this system: uniform electric field fluctuations.

The CQ qubit is formed in three adjacent semiconducting quantum dots sharing a single electron [21]. In the localized charge basis \([|100\rangle, |010\rangle, |001\rangle]\), where the basis states denote the electron being in the first, second, or third dot, respectively, the Hamiltonian is given by

\[ H_{\text{CQ}} = \begin{pmatrix} \epsilon_d & t_A & 0 \\ t_A & \epsilon_q & -t_B \\ 0 & t_B & -\epsilon_d \end{pmatrix} + \frac{U_1 + U_3}{2}, \]

where \( U_{1,2,3} \) are the on-site potentials for the three dots, \( t_{A,B} \) are tunnel couplings between adjacent dots, and \( \epsilon_d = (U_1 - U_3)/2 \) and \( \epsilon_q = U_2 - (U_1 + U_3)/2 \) denote the dipolar and quadrupolar detuning parameters, respectively. We now define a new set of basis states [21]

\[ |C\rangle = |010\rangle, \quad |E\rangle = \frac{|100\rangle + |001\rangle}{\sqrt{2}}, \quad |L\rangle = \frac{|100\rangle - |001\rangle}{\sqrt{2}}, \]

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and $\delta \epsilon$ point to a region with $\epsilon_0$ as our model Hamiltonian. We restrict our analysis to the uniform electric field fluctuations [21]. Fluctuations of the form a decoherence-free subspace (DFS) with respect to where $|\epsilon_0 = g = 0\rangle$ for the CQ qubit. The insets (which use the $\delta \epsilon_d/h$ corresponding fluctuations. In Ref. [21] it is argued that, while $\delta \epsilon$ can now be performed by pulsing gate operations. A rotation about the $z$ axis of the Bloch sphere generates $|\epsilon_0 = g = 0\rangle$ for the CQ qubit. We note that $\delta \epsilon_d$ corresponds to logical states and $|L\rangle$. The unitary operators for noisy $z$ and $x$ rotations are $U_z(\epsilon_d, \delta \epsilon_d, \varphi) = e^{-i(\epsilon_d + \delta \epsilon_d g)\varphi/\epsilon_d}$ and $U_x(g, \delta \epsilon_d, \theta) = e^{-i(\epsilon_d + \delta \epsilon_d g)\theta/2g}$, for arbitrary rotation angles $\varphi$ and $\theta$. For bang-bang gates, these angles are related to the corresponding gate times as $\varphi = t_z(\epsilon_d/g)$ and $\theta = t_x(2g/h)$. Hence, $\varphi$ must have the same sign as $\epsilon_d$, and $\theta$ must be positive, since $g > 0$ for most solid-state devices, including quantum dots. Our approach will be to compose $U_z$ and $U_x$ gates to obtain arbitrary single-qubit rotations in the logical subspace as a function of $\epsilon_d, g, \varphi$, and $\theta$, and then determine the conditions that should be imposed on these parameters to suppress the leading order leakage terms in the composed gates.

A schematic diagram of a bang-bang pulse sequence for this $R_{zxz}$ operation is shown in the insets of Fig. 1. We now expand $R_{zxz}$ about $\delta \epsilon_d = 0$ and note that the first order terms of $\delta \epsilon_d$ vanish from the Taylor expansion if we set $U_z(\epsilon_d, \delta \epsilon_d, \varphi/2)U_x(g, \delta \epsilon_d, \theta)U_z(-\epsilon_d, \delta \epsilon_d, -\varphi/2)$. A schematic diagram of a bang-bang pulse sequence for this $R_{zxz}$ operation is shown in the insets of Fig. 1.

We now expand $R_{zxz}$ about $\delta \epsilon_d = 0$ and note that the first order terms of $\delta \epsilon_d$ vanish from the Taylor expansion if we set $\epsilon_q = -g\varphi/2 \cot \left( \frac{\theta}{4} \right)$. This condition, along with the conditions on $\varphi, \theta$, and $g$, are all satisfied when $2\pi < \theta < 4\pi$, which means that $U_z$ rotations are between 1 and 2 cycles around the Bloch sphere.

We note that $R_{zxz}(\theta, \varphi)$ corresponds to a rotation of angle $\theta$ about an arbitrary axis $\hat{n} = \cos(\varphi/2)\hat{x} + \sin(\varphi/2)\hat{y}$ in the $x$-$y$ plane of the logical subspace (modulo an irrelevant overall phase). To estimate analytically how the computational and leakage errors depend on the noise $\delta \epsilon_d/g$, we evaluate the higher order terms in the Taylor expansion of $R_{zxz}(\theta, \varphi)$ [28],
observing that the error probability within the logical subspace $\langle R_{\text{le}}(\delta_\epsilon) \neq 0 \rangle - \langle R_{\text{le}}(\delta_\epsilon = 0) \rangle^2$ scales as $(\delta_\epsilon/g)^4$, while the probability of leakage errors $P_{\text{le}} = \langle [L,R_{\text{le}}(C,E)]^2 \rangle$ scales as $(\delta_\epsilon/g)^6$. Therefore, for a typical experimental value of $\delta_\epsilon/g \sim 10^{-1}$, the error in the computational subspace is $\sim 10^{-4}$ and $P_{\text{le}} \sim 10^{-6}$. This represents a remarkable improvement over the bare gates, for which the leakage error is of order $10^{-2}$ when $\delta_\epsilon/g \sim 10^{-1}$ (for a detailed comparison of leakage errors in simple and composite pulse sequences, see Ref. [28]).

Finally, we construct a completely general and arbitrary rotation on the Bloch sphere. Since the effective rotation axis for an $R_{\text{le}}$ gate lies anywhere in the $x$-$y$ plane, an arbitrary rotation requires just two steps [29]: $R_{\text{le}}(\phi_1, \phi_2)R_{\text{le}}(\phi_0, \phi_1)$, comprising six bare gates. Although there are penalties associated with longer gate sequences, such as reduced gate speeds and accumulation of errors, the sequences described here improve the overall scaling of errors with respect to the noise amplitude. Hence, they are always beneficial if $\delta_\epsilon g$ is small enough.

While an $R_{\text{le}}$ gate produces arbitrary rotations in the $x$-$y$ plane, Eq. (7) indicates that the control parameter $g$ ($\epsilon_0$) diverges in the limit $\theta \rightarrow 2\pi$ (4$\pi$). This implies that an identity gate cannot be implemented using a single $R_{\text{le}}$ operation. A high-fidelity identity gate is an essential ingredient for fault-tolerant quantum computing because most qubits remain idle during most of the quantum error-correction cycle. We cannot use the “idle gate”, corresponding to $\epsilon_0 = g = 0$, as an identity gate because $\delta_\epsilon \neq 0$ causes errors. It is possible to obtain an identity gate from the sequence $\langle R_{\text{le}}(\pi,\phi) \rangle^2$, which comprises six bare gates; however it is interesting to ask whether shorter identity sequences exist.

We now present a viable identity sequence $R_I$ that uses just four bare gates: Following the same procedure used to construct $R_{\text{le}}$, we consider the sequence

$$R_I = \left[ U_g(\epsilon_0, \delta_\epsilon, 2\pi) U_g(\delta_\epsilon, \delta_\epsilon, 2\pi) \right]^2. \quad (8)$$

$R_I$ is clearly an identity operation when $\delta_\epsilon = 0$. To address the case $\delta_\epsilon \neq 0$, we perform a Taylor expansion of $R_I$ about $\delta_\epsilon = 0$ [28], noting that the first-order term in $\delta_\epsilon$ vanishes due to the special form of the sequence, even without imposing special constraints on $\epsilon_0$ and $g$. The Taylor expansion also indicates that the probability of computational errors, given by $\langle [C(R_I)]^2 \rangle = \langle [E(R_I)]^2 \rangle$, scales as $(\delta_\epsilon/g)^4$, while the probability of leakage errors, given by $\langle [L(R_I)]^2 \rangle = \langle [L(R_I)]^2 \rangle$, scales as $(\delta_\epsilon/g)^6$.

To understand why our pulse sequences work, we now describe a method for simultaneously visualizing the evolution of the logical and leakage states on a pair of coupled Bloch spheres. An arbitrary state $|\Psi\rangle$ of a three-level quantum system can be expressed in the basis $\{|C\rangle, |E\rangle, |L\rangle\}$ using four angle variables:

$$|\Psi\rangle = \begin{pmatrix} \cos(\theta/2) \cos(\chi/2) \\ \cos(\theta/2) \sin(\chi/2) e^{i\epsilon} \\ \sin(\theta/2) e^{i\zeta} \end{pmatrix}. \quad (9)$$

When $\theta = 0$, $|\Psi\rangle$ maps onto $|\Psi_g\rangle = \cos(\chi/2)|C\rangle + \sin(\chi/2)e^{i\zeta}|E\rangle$, comprising a two-level system. Similarly, when $\chi = 0$, $|\Psi\rangle$ maps onto $|\Psi_r\rangle = \cos(\theta/2)|C\rangle + \sin(\theta/2)e^{i\zeta}|L\rangle$, comprising a different two-level system. The full mapping $|\Psi\rangle \rightarrow \{|\Psi_g\rangle, |\Psi_r\rangle\}$ is bijective, and the states $|\Psi_g\rangle$ and $|\Psi_r\rangle$ may alternately be represented as unit vectors on a pair of Bloch spheres that we label ‘green’ and ‘red’, respectively. The angles $\theta$, $\chi$, $\epsilon$, and $\zeta$ here correspond to the usual polar and azimuthal angles of the respective Bloch spheres.

In the absence of leakage, the state vector of the red sphere points to the north pole, and the state vector of the green sphere describes the logical qubit in the usual way. When leakage occurs, the state vector of the red sphere deviates from north, with its latitude describing the amplitude of the leakage state and its azimuth describing the relative phase of $|L\rangle$ with respect to $|C\rangle$. The green sphere still describes the logical qubit, with its state vector renormalized to lie on the unit sphere. For quantum systems with multiple leakage states, we could extend this geometrical representation to include multiple Bloch spheres, with one sphere describing the logical states and one additional sphere for each leakage state.

We restrict our analysis to a single leakage state and compute the time-dependent trajectories of the green and red Bloch spheres for the $R_{\text{le}}$ pulse sequence.

Our analytical results for pulse sequences were previously obtained in the bang-bang limit. However in charge qubit experiments, the combination of high-frequency filtering and short gate times (~0.1–1 ns) means that pulses will experience significant rounding. To account for this, we now consider pulse sequences with smooth profiles and finite rise times. Under such conditions, constraints on the pulse shape, such as Eq. (7), must be modified. To do this, we keep our previous sequences but vary the control parameters and gate times. The parameters are chosen to maximize the process fidelity, defined as [30,31]

$$F = \frac{\text{Tr}(U_U^{\dagger}U_U)}{d(d+1)}. \quad (10)$$

Here, $U_U$ is the desired gate operation in the 2D logical subspace, $U$ is the unitary operation obtained from simulations and projected onto the 2D logical subspace, and $d = 2$. Such projections do not generally yield unitary operators, however Eq. (10) takes this into account and captures both computational and leakage errors.

A typical result from our gate optimization procedure is shown in Fig. 2. Here, we consider an effective $X[-\pi/2]$ rotation, obtained from the three-step sequence defined in Fig. 2(a), which is analogous to the bang-bang sequence $R_{\text{le}}(2\pi \times \pi/2, 2\pi \times \pi/2)$. We adopt several fixed device parameters, including error-function pulse profiles with rise times of 50 ps [28], a peak tunnel coupling of $g/h = 3$ GHz, and a typical experimental noise value of $\delta_\epsilon g/h = 0.3$ GHz [26]. We then numerically optimize the gate times $t_c$ and $t_s$ defined in the figure, and the peak value of $\epsilon_0$, to maximize the fidelity. Figures 2(b) and 2(c) show the time evolution of the resulting gate for the initial state $\cos(\pi/10)|C\rangle + \sin(\pi/10)|E\rangle$. Although there is significant leakage in the middle of the operation, the final leakage probability is strongly suppressed. From Eq. (10), the final infidelity is found to be $1 - F = 2.46 \times 10^{-8}$, while the total leakage errors corresponding to $|C\rangle$ and $|E\rangle$ states are $1.5 \times 10^{-10}$ and $1.2 \times 10^{-10}$, respectively. These results are obtained for the smooth pulse sequence.
Leak prob.

 FIG. 2. A numerically optimized gate operation. (a) The pulse sequence for an \( X_{\pi/2} \) gate, assuming rounded pulse profiles with 50 ps rise times. Here, the optimization parameters are the gate times \( t_1 \) and \( t_2 \), and the peak value of \( \epsilon_q \). The peak value of \( g/h = 3 \text{ GHz} \) is fixed. (b) Leakage probabilities as a function of time. The final leakage error is of order \( 10^{-10} \). (c) Time evolution of the logical and leakage states on green and red “Bloch spheres”, respectively, for the smooth pulse shown in (a). Here, the color of the state vectors indicates the time. The transitions between \( z, x, \) and \( z \) rotations are indicated by symbols, as consistent with (a).

shown in Fig. 2(a) and represent orders of magnitude of improvement in the infidelity compared to the bang-bang pulse under the constraint (7). This improvement can be attributed to the numerical optimization procedure. Indeed, by relaxing constraint (7), while retaining the three-pulse construction, the dynamics can be modified to cancel out higher order noise effects for either bang-bang or smooth pulses.

The benefit of the two-Bloch-sphere representation is apparent in Fig. 2(c), where it provides a qualitative understanding of how our pulse sequence suppresses leakage errors. Starting from the north pole of the red sphere, the leakage probability grows slowly during the first \( z \) rotation. As it grows, the leakage state accumulates phase, causing the state vector to spiral outward and downward. The special constraint on the control parameters, analogous to Eq. (7), causes the leakage phase to be inverted during the \( x \) rotation, similar to a Hahn echo. The final \( z \) rotation then yields a time evolution that is opposite of the initial evolution (a reverse spiral), which eventually passes very close to the north pole.

In conclusion, we have developed special pulse sequences for a three-level system, comprised of two logical states and a leakage state, which suppress the leakage probability and enable arbitrary, high-fidelity, single-qubit rotations. We consider a general model, in which only one of the logical states is coupled to the leakage state. A crucial feature of this method is that it requires no knowledge of the coupling to the leakage state. The method is perfectly suited for the quantum dot charge quadrupole (CQ) qubit, whose leakage coupling is caused by charge noise. Taking the CQ as our model system we have shown that computational errors in the gate operation arise at fourth order in the noise amplitude, while leakage errors arise at sixth order. This represents a substantial improvement over conventional pulse schemes, for which errors arise at second order in the noise amplitude. While it is possible to design high-fidelity quantum gates for quantum dot qubits using optimal control schemes [16,17], we emphasize that our approach yields high gate fidelities without knowing the amplitude of the noise, which sets it apart from those schemes. We also note that a Hahn echo sequence, which has the same level of complexity as our sequence, has recently been employed successfully to improve the fidelities of charge dipole qubits [32]. This suggests that our scheme could also yield high-fidelity gate operations in charge quadrupole qubits. To understand why our pulse sequence works, we have introduced an intuitive geometrical mapping of the three-level system, in which the logical and leakage states are represented on two coupled Bloch spheres. In this way, the cancellation of leakage errors appears similar to a Hahn echo. Our proposal therefore provides an important step towards fault-tolerant quantum computation in quantum dots.

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